

Self assembling graphs

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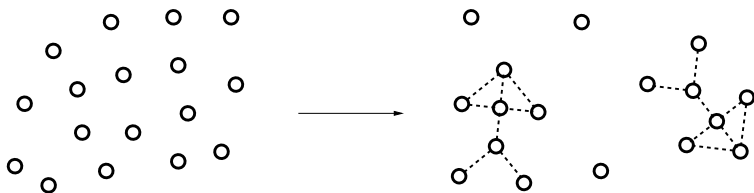
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The problem

- ▶ **Question** : How a collective behaviour may emerge from elementary interactions (forward and backward)
- ▶ **Applications**
 - ▶ Molecular biology (backward)
 - ▶ Genetic engineering (forward)
 - ▶ Distributed robotics (forward)

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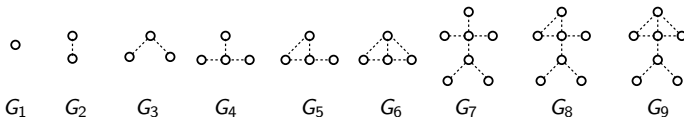
Formal approach

- ▶ Model of connections, space, ...
What objects are built ?
- ▶ Combinatorics of the components:
What can compute a basic element ?
Reasonable assumption depends on the context.
- ▶ Model of communication:
How interact an agent ?
How information is transmitted ?

We look for a syntaxe able to describe the **elements** and the **assembling**.

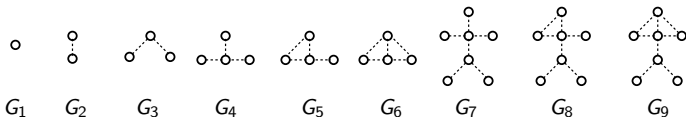
Preliminary work

► \mathcal{G} : Set of **explorative graphs** :

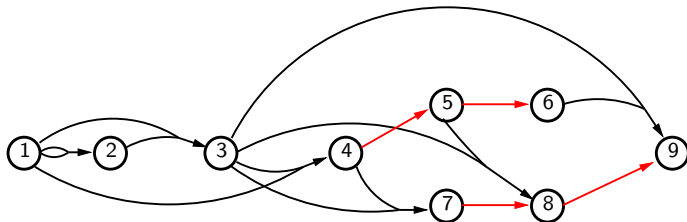


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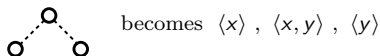
- ▶ **Assembling graph** of the final target:



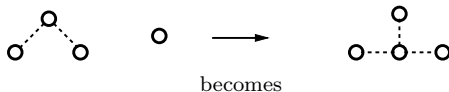
The syntax

Syntactic representation of graphs :

- ▶ Nodes = agents
- ▶ Edges = private names sharing



Construction rules :



$$\langle x \rangle, \langle x, y \rangle, \langle y \rangle, \langle \rangle \longrightarrow (\nu z) (\langle x \rangle, \langle x, y, z \rangle, \langle y \rangle, \langle z \rangle)$$

Formalisation of the problem

- ▶ Extraction of a **core language**:
 $\langle x \rangle, \langle x \rangle, \langle \rangle \rightarrow (\nu y)(\langle x \rangle, \langle x, y \rangle, \langle y \rangle)$
 \implies restriction on synchronisation ability
- ▶ Expected property: equivalent behaviour

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- ▶ Comparison of states

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- \implies Mathematical tool: **bisimulation**

Bisimulation

[*Observation*]

The least relation \downarrow such that $S \downarrow C$ only if the connected component C appears in the system S

[*(weak) bisimulation*]

Binary symmetric relation \mathfrak{R} such that if $S \mathfrak{R} S'$ then

1. If $S \rightarrow T$ then $\exists T'$ s.t. $S' \rightarrow^* T'$ and $T \mathfrak{R} T'$
2. If $S \downarrow C$ then $S' \downarrow C$.

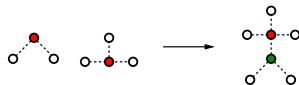
Intuitive features of the algorithm

- ▶ Only one active agent by component.
- ▶ Local knowledge of the component's structure.
- ▶ Each agent knows its role in the component.
- ▶ Propagation of the changes related to an interaction by the use of a spanning tree.

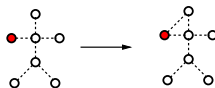
Traduction of the reactions

Set of reactions :

- ▶ Connection between 2 disjoint complexes



- ▶ Cyclic connection



- ▶ Propagation updates
- ▶ Activity switch
- ▶ Mechanism to handle the deadlocks

Connection rule

$\forall G_1, r_1, G_2, r_2$ such that $[G_1.r_1 \oplus G_2.r_2]^{abs} \in \mathcal{G}$

$$\begin{array}{c}
 \langle S_1, C_1, g_1, r_1, Act(G_1) \rangle \\
 \langle S_2, C_2, g_2, r_2, Act(G_2) \rangle \\
 \downarrow \\
 (\nu com) \\
 \langle S_1 \cup \{com\}, C_1, g_1, r_1, Act(G_1.r_1 \oplus G_2.r_2) \rangle \\
 \langle S_2 \cup \{com\}, C_2, g_1, r_2 + \|G_1\|, Up(S_2, \|G_1\|) \rangle
 \end{array}$$

Self connection

$\forall G, r_1, r_2$ such that $[r_1 \overset{G}{\wedge} r_2]^{abs} \in \mathcal{G}$

$$\begin{array}{l} \langle S_1, C_1, g, r_1, Act(G) \rangle \\ \langle S_2, C_2, g, r_2, P \rangle \end{array}$$

↓

(νcom)

$$\begin{array}{l} \langle S_1, C_1 \cup \{com\}, g, r_1, Act(r_1 \overset{G}{\wedge} r_2) \rangle \\ \langle S_2, C_2 \cup \{com\}, g, r_2, P \rangle \end{array}$$

Updates

$$\begin{array}{l}
 \langle S_1, C_1, g_1, r_1, Up(L + \{com\}, Num) \rangle \\
 \langle S_2 + \{com\}, C_2, g_2, r_2, P \rangle \\
 \downarrow \\
 \langle S_1, C_1, g_1, r_1, Up(L, Num) \rangle \\
 \langle S_2 + \{com\}, C_2, g_1, r_2 + Num, Up(S_2, Num) \rangle
 \end{array}$$

End of updates

$$\langle S, C, g, r, Up(\emptyset, Num) \rangle \longrightarrow \langle S, C, g, r, P \rangle$$

Activity switch

$$\begin{array}{l} \langle S_1, C_1, g, r_1, Act(G) \rangle \\ \langle S_2, C_2, g, r_2, P \rangle \\ \downarrow \\ \langle S_1, C_1, g, r_1, P \rangle \\ \langle S_2, C_2, g, r_2, Act(G) \rangle \end{array}$$

A more intuitive presentation

Conclusions

- ▶ Formal language to express the problem of self assembling parts.
- ▶ Mathematical tool to resolve it.
- ▶ Similar work in \mathbb{R}^n (to be done as in the demo).
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